

```

-- literate haskell .lhs
-- radky zacinaji >
{-
> lhsFunkce x = lhsTelo (x * x)
>   where
>     lhsTelo y = y `div` 2

-- z prelude
foldr :: (a->b->b) -> b -> [a] -> b
foldr f v []      = v
foldr f v (x:xs) = f x (foldr f v xs)

(.) :: (b->c) -> (a->b) -> (a->c)
(f . g) x = f (g x)
-- > test0 x = faze3 ( faze2 ( faze1 x ) )
-- > test1 x = (faze3 . faze2 . faze1) x
-- > test2    = faze3 . faze2 . faze1 -- stejne typy

infixr 0 $
($) :: (a->b) -> a -> b
f $ x = f x
-- > test3 x = faze3 $ faze2 $ faze1 x

infixr 0 $!
($!) :: (a->b) -> a -> b
f $! x = ... interni, pro striktni vyhodnocovani
-- vyhodnoti top konstruktor x
-}

```

```

-----  

-- 11.5.2009 stromy  

-- (1) data len v listoch  

data LTree a = LLeaf a  

    | LBranch (LTree a) (LTree a)  
  

-- (2) data interne, BVS  

data BTree a = BLeaf  

    | BBranch (BTree a) a (BTree a)  

-- BVS = BTree (TKey,TValue) -- dvojice :-(  
  

-- (3) data: ine v listoch, ine vnutri  

data DTree a b = DLeaf a  

    | DBranch (DTree a b) b (DTree a b)  

binVyraz :: DTree Int String  

binVyraz = DBranch (DLeaf 1)  

    "+"  

    (DBranch (DLeaf 2)"*" (DLeaf 3))  
  

-- R-B trees ...  

t1BTree = BBranch  

    (BBranch BLeaf 1 BLeaf)  

    2  

    (BBranch  

        (BBranch BLeaf 3BLeaf)  

        4  

        (BLeaf) )  
  

{- typovy konstruktor BTree, LTree druhu(::) * -> *  

typ.k. DTree druhu (kind) *->*>->*  

-- datovy konstruktor BLeaf :: BTree a -- polymorfni  

dat.k. BBranch :: BTree a -> a -> BTree a -> BTree a  

dat.k. DLeaf :: a -> DTree a b  

-}

```

```

foldBT :: b->(b->a->b->b)->BTree a-> b
-- ala foldr pro seznamy (strukt. rek.)
foldBT fLeaf fBranch bt = f bt where
  f BLeaf          = fLeaf
  f (BBranch l x r) = fBranch (f l) x (f r)

copyBT t = foldBT BLeaf BBranch t -- "identita"
sizeBT t = foldBT 0 (\l x r->l+1+r) t
depthBT t = foldBT 1 (\l x r->1+max l r) t
listify t = foldBT [] (\l x r->l++(x:r)) t -- infixBT
postfix t = foldBT [] (\l x r->l++r++[x]) t
sumBT t  = foldBT 0 (\l x r->l+x+r) t
mapBT f t  = foldBT BLeaf (\l x r->BBranch l (f x) r) t
heapify t = foldBT BLeaf (\l x r->inHeap l x r) t
-- where inHeap na dalsi strane
prumer t = fromInteger suma/fromInteger size
-- obecne: potrebuj i postspracovani
where
  (suma,size) = foldBT (0,0)
    (\(ls,lv) x (rs,rv)->(ls+x+rs,lv+1+rv)) t
--bvs t = foldBT True (\l x r->l && r && locOrdered l x r) t
-- where locOrdered l x r = ...
vaha :: BTree (Int,a) -> (Int,Int)
  -- (freq,key)-> (freq, vaha); vaha = sum(freq*depth)
vaha t = foldBT (0,0) (\(fl,vl)(fx,_k)(fr,vr)
  ->(fl+fx+fr,vl+fl+fx+vr+fr) ) t
-- ... a varianty fold:
-- Obecne musim vracet i puvodni i spracovanou strukt. (tj. dvo

```

```

-- mkHeap :: BTree a -> BTree a
-- jen preusporadani, ne stavba
mkHeap BLeaf = BLeaf
mkHeap (BBranch l x r) = inHeap (mkHeap l) x (mkHeap r)

inHeap BLeaf x BLeaf = BBranch BLeaf x BLeaf
inHeap l@(BBranch l1 xl l2) x BLeaf
| x <= xl = BBranch l x BLeaf
| True     = BBranch (inHeap l1 x l2) xl BLeaf
inHeap BLeaf x r@(BBranch r1 xr r2)
| x <= xr = BBranch BLeaf x r
| True     = BBranch BLeaf xr (inHeap r1 x r2)
inHeap l@(BBranch l1 xl l2) x r@(BBranch r1 xr r2)
| x <= xl && x <= xr
    = BBranch l x r
| xl <= xr = BBranch (inHeap l1 x l2) xl r
| True     = BBranch l xr (inHeap r1 x r2)

-- aktivni konstruktory: buduji pozmenenou strukturu
t1Heap =
  bBranch
    (bBranch bLeaf 1 bLeaf)
  2
  (bBranch
    (bBranch bLeaf 3 bLeaf)
    4
    (bLeaf)
  )
where
bLeaf = BLeaf
bBranch = inHeap

```

```
-----  
-- stromy n-arne (Rose Tree)  
data RTree a = RT a [RTree a]  
  
-- aplikace: XML...  
data XML a = Elm a [XML a]  
           | Txt String  
-- a/String -- jen tag  
-- a/(String,[(String,String)])  
-- show bez chyb, praca so zoznamami elementov  
  
{- -- aplikace: univ.vyrazy, synt. stromy ...  
data Expr a = Var String  
           | Fn String [Expr a]  
           ...  
           | Con a  
           | If (Expr Bool) (Expr a) (Expr a)  
           | Inc (Expr Int)  
-- vyrazy v FP, nejen PP  
-- fantomove typy  
-}
```

```

-- huffmanovo kodovani: optimalni, hladovy alg.
-- vzdaleny ;-) cil
{- test_HDE s = let t = mkHTree $ freq s
               in s == (dec t $ enc (hstrom2tab t) s)
-}

data HStrom a = HL a
               | HV (HStrom a) (HStrom a)

hstrom2tab :: HStrom (Char {-,Int-}) -> [(Char, String)]
hstrom2tab (HL(c)) = [(c, "")]

hstrom2tab (HV l r) =
  map (\(c,kod)->(c,'0':kod))
    (hstrom2tab l)
  ++
  map (\(c,kod)->(c,'1':kod))
    (hstrom2tab r)

-- volani> hs2t "" strom -- akumulator
hs2t kod (HL(c,"")) = [(c,kod)]
hs2t kod (HV l r) =
  hs2t (kod++"0") l ++
  hs2t (kod++"1") r

enc :: (Eq a, Show a) => [(a, [b])] -> [a] -> [b]
enc tab []      = []
enc tab (x:xs) =
  enc1 tab x ++
  enc tab xs
  where
    enc1 []           x =
      error("enc: missing code for:"++show x)
    enc1 ((x1,kod):tab) x =
      if x==x1 then kod
      else enc1 tab x

```

```
enc2 tab xs =
    concat(map(\x -> snd(head(filter (f x) tab))) xs)
  where f x = \(x1,_kod) -> x==x1 -- x: lambda-lifting
        -- f x = (x==).fst -- varianta
-- = concat $ map (\x -> snd $ head $ filter f tab) xs

-- concat []      =  [] -- v Prelude
-- concat (x:xs) =  x ++ concat xs
-- anebo: concat XSS = foldr (++) [] XSS
```

```

dec :: HStrom a -> [Char] -> [a]
dec strom []      = []
dec strom inp    = c : dec strom inp1
where
  (c, inp1)      = dec1 strom inp
  dec1 :: HStrom a -> [Char] -> (a, [Char])
  dec1 (HL c) inp          = (c,inp)
  dec1 (HV l r) ('0':inp) = dec1 l inp
  dec1 (HV l r) ('1':inp) = dec1 r inp
  dec1 (HV l r) ( x :inp) = error ("dec: bad code:"++show x)
  dec1 (HV l r) []         = error "dec: unexpected EOF"

dec2 :: HStrom a -> [Char] -> [a]
dec2 strom inp = case x of
  Nothing          -> []
  -- interni chyba: nic se nedekoduje
  Just (c,inp1)   -> c : dec strom inp1
where
  x = dec1 strom inp
  dec1 :: HStrom a -> [Char] -> Maybe (a, [Char])
  dec1 (HL c) inp          = Just (c,inp)
  dec1 (HV l r) ('0':inp) = dec1 l inp
  dec1 (HV l r) ('1':inp) = dec1 r inp
  dec1 (HV l r) ( x :inp) = Nothing
  dec1 (HV l r) []         = Nothing
-- pozn: odebirana cast je dana kontextem
{-
data Maybe a = Nothing
  | Just a
  deriving (Eq, Show)
-}

```

```

mkHTree :: [(Char,Int)] -> HStrom Char
mkHTree xs = faze5
where
  cmp (c1,f1) (c2,f2) = f1<=f2
  -- cmp    = \((c1,f1) (c2,f2) -> f1<=f2
  faze1 = sort cmp xs
  faze2 = map (\(c,f) -> (HL c, f)) faze1
  faze3 = iterate spoj faze2
  faze4 = dropWhile (\x -> length x > 1) faze3
  faze5 = fst $ head $ head faze4
  spoj [x]           = [x]
  spoj ((c1,f1):(c2,f2):xs) = insert cmp (HV c1 c2,f1+f2) xs
  sort cc xs = foldr (insert cc) [] xs
  insert cc x [] = [x]
  insert cc x vs@(y:ys)
    | x `cc` y = (x:vs)
    | True = y:insert cc x ys

```

```

freq :: String -> [(Char,Int)] -- histogram
freq xs = foldr ordBagUnion [] $ map (\x->[(x,1)]) xs
-- O(n^2) :-(

ordBagUnion [] ys = ys
ordBagUnion xs [] = xs
ordBagUnion xx@(x,fx):xs yy@(y,fy):ys
| x < y = (x,fx) : ordBagUnion xs yy
| x == y = (x,fx+fy) : ordBagUnion xs ys -- memory leak
| x > y = (y,fy) : ordBagUnion xx ys

foldDvoj f e [] = []
foldDvoj f e xs =
  (head.head.dropWhile((1<).length).iterate(spojDvoj f))xs
where
  spojDvoj f (x1:x2:xs) = f x1 x2 : spojDvoj f xs
  spojDvoj f xs          = xs
freq2 xs = foldDvoj ordBagUnion [] $ map (\x->[(x,1)]) xs
testFreq k c = [s | s <- var0pak1 k ['a'..c],
                  let s1=s,
                  freq2 s /= freq s ]
-- testFreq 7 'f' : 1'23, length 279936 testov
{- prelude
iterate :: (a->a) -> a -> [a]
iterate f x = x : iterate f (f x)

dropWhile :: (a->Bool) -> [a] -> [a]
dropWhile _ []      = []
dropWhile p (x:xs)
| p x              = dropWhile p xs
| otherwise         = x : xs
-- DC: ukoncenie podla vztahu dvoch prvkov
-}

```

```

varOpak :: [a] -> [b] -> [[[a,b]]]
varOpak [] _ = [[]] -- vhodne "dodefinovano"
varOpak (x:xs) r =
  [(x,y):vs | y<-r, vs <- varOpak xs r]

```

```

varOpak1 :: Int -> [b] -> [[b]]
varOpak1 k r =
  map           -- pro všechny variace
    (map       -- pro všechny prvky ve var.
      snd)     -- zahazuju dom., vracím hodnotu
  (varOpak     -- volání puv. fce dostane:
    [1..k]   -- domain jako seznam
    r)       -- range bez zmen

```

```

komb :: Int -> [a] -> [[a]]
komb 0 _ = [[]] -- vhodne "dodef."
komb (n+1) [] = []
komb (n+1) (x:xs) =
  [x:ko | ko <- komb n xs] ++ komb (n+1) xs

```

```

kombOpak :: Int -> [a] -> [[a]]
kombOpak 0 _ = [[]] -- vhodne "dodef."
kombOpak (n+1) [] = []
kombOpak (n+1) (xs) =
  [head xs:ko | ko <- kombOpak n xs] ++ kombOpak (n+1) (tail xs)

```

```

-- testy: QuickCheck 2000, SmallCheck 2008, ...
test_HDE s = let t=mkHTree $ freq s
             in s == (dec t $ enc (hstrom2tab t) s)
prop_HDE k c = [s|s<-kombOpak k ['a'..c], not $ test_HDE s]
               -- varOpak1

-- > prop_HDE 3 'b'
-- >> ["aaa","bbb"]

-- s kombOpak: 15 'g' : 35'', 54264 testov
-- s varOpak1: 7 'f' : okolo 2', 279936 testov
{-
-- Zde ukonceno na Prednasce 12.05.2009,
-- nektere drobnosti preskocene
Obsahlejsi popis napr. Huffmanova kodovani v textu "Haskell v pr
----- -}

```

```
-----  
stableSort cmp xs =  
    map fst $ -- odstraneni poradi (A)  
        sort1 cmp2 $ -- trideni se zmenou porovnavaci fce (B)  
            zip xs [1..] --pridani poradi (C)  
where cmp2 (x,i) (y,j) = x `cmp` y && not (y `cmp` x)  
                           || x `cmp` y && y `cmp` x && i <= j  
sort1 cmp []      = []  
sort1 cmp (p:ys) =  
    sort1 cmp ([y | y <- ys, y `cmp` p] ++ (p : [y | y <- ys, not$y `cmp` p]))
```

{- -- add
F# (.NET), Scala (nad JVM)
lift: na Maybe, Either
na zoznamy (vektory), matice
na funkcie - casova zavislost
DSEL - Domain Specific (Embedded) Language
fantomove typy
Calling hell from heaven and heaven from hell, 1999
kontext
Blub: hypoteticky prog. jazyk, Paul Graham
-}

-- (tvorba indexu: reverzny zoznam)

type Word = String

splitWords :: String -> [Word]

splitWords st = split (dropSpace st)

split :: String -> [Word] - dostává ř. bez úvodních mezer

split [] = []

split st = (getWord st):split(dropSpace(dropWord st))

dropSpace st = dropWhile (\x->elem x whitespace) st

dropWord st = dropWhile (\x->not(elem x whitespace)) st

dropWhile :: (t -> Bool) -> [t] -> [t] - obecná výkonná procedura

dropWhile p [] = [] -- zahodí poč. prvky, které nesplňují podm.

dropWhile p (a:x) -- v prelude

| p a = dropWhile p x

| otherwise = (a:x)

getWord st = takeWhile (\x-> not(elem x whitespace)) st

-- DC: takeWhile p x = -- v prelude

whitespace = [, \t, \n]

Aritmetika s testováním chyb (pomocí Maybe)

```
lift2M :: (a->b->c)->Maybe a -> Maybe b -> Maybe c
lift2M op Nothing _           = Nothing chyba v 1. arg.
lift2M op _      Nothing    = Nothing chyba v 2. arg.
lift2M op (Just x) (Just y) = Just (x op y) bez chyb
```

```
minusM :: Maybe Float -> Maybe Float -> Maybe Float
minusM x y = lift2M (-) x y
```

-- Vyvolání chyby

```
delenoM x y = if y==Just 0 then Nothing --test a vyvolání chyby
                 else lift2M (/) x y
? delenoM(Just 3)(minusM(Just 2)(Just 2)) -- 3/(2-2)
? delenoM(Just 3)(lift2M(-)(Just 2)(Just 2)) -- dtto Nothing
```

```
data AV a = AV a :-: AV a | AV a :/: AV a | Con a |
```

```
eval :: AV Integer -> Maybe Integer
eval (av1 :/: av2) = delenoM (eval av1) (eval av2)
eval (av1 :-: av2) = lift2M (-) (eval av1) (eval av2)
eval (Con x)       = Just x -- cena za Maybe: zábalení výsledku
```

-- Odchycení chyby:

```
? catch (eval (Con 3 :/: (Con 2 :-: Con 2))) 1
catch Nothing oprData = opravnaFce oprData
catch (Just x) _      = x
```

map pro jiné d.s.

pro binární stromy

rozbor podle konstruktoru

mapTree :: ($a \rightarrow b$) \rightarrow Tree a \rightarrow Tree b

mapTree f (Leaf a) = Leaf (f a)

mapTree f (Branch l r) = Branch (mapTree f l)
(mapTree f r)

n-ární stromy

data NTree a = Tr a [NTree a]

mapNT :: ($a \rightarrow b$) \rightarrow NTree a \rightarrow NTree b

mapNT f (Tr x trees) = Tr (f x) (map (mapNT f) trees)

typ Maybe, Union - nerekurzivní

mapM :: ($a \rightarrow b$) \rightarrow Maybe a \rightarrow Maybe b

mapM f Nothing = Nothing

mapM f (Just x) = Just (f x)

mapU :: ($a \rightarrow c$) \rightarrow ($b \rightarrow d$) \rightarrow Union a b \rightarrow Union c d

mapU f g (Left x) = Left (f x)

mapU f g (Right y) = Right(g y)

Stavové programování

(Návrhový vzor iterátor)

N.vzory ve FP lze (často) napsat jako kod s funkč. parametry
vs. (často) pseudokód v OOP

`iterator::s->(s->s)->(s->Bool)->s`

-- vrací celý (interní) stav, zde není výstupní projekce

`iterator init next done =`

`head(dropWhile (not.done) --mezivýsl. se průběžně`

`(iterate next init)) -- zahazují`

DC: `fixpoint :: (s->s)->s->s :vrací v=f^n(x) pro min. n:v=f(v)`

Počítání s programy

Typicky: dokazování vlastností programů

(Částečná) správnost vzhledem k specifikaci

Transformace programů pro optimalizaci

př.: asociativita (++) append

$$[] \text{ ++ } ys = ys$$

$$(x:xs) \text{ ++ } ys = x: (xs \text{ ++ } ys)$$

Tvrzení: $x \text{ ++ } (y \text{ ++ } z) = (x \text{ ++ } y) \text{ ++ } z$ pro konečné x, y, z .

$$x = [] : LS = [] \text{ ++ } (y \text{ ++ } z) = y \text{ ++ } z$$

$$PS = ([] \text{ ++ } y) \text{ ++ } z = y \text{ ++ } z$$

$$x = a:v : LS = (a:v) \text{ ++ } (y \text{ ++ } z) = / \text{ def. } ++$$

$$a:(v \text{ ++ } (y \text{ ++ } z)) = / \text{ ind. předp. } ++$$

$$a:((v \text{ ++ } y) \text{ ++ } z)$$

$$PS = ((a:v) \text{ ++ } y) \text{ ++ } z = / \text{ def. } ++$$

$$((a:(v \text{ ++ } y)) \text{ ++ } z) = / \text{ def. } ++$$

$$a:((v \text{ ++ } y) \text{ ++ } z) \text{ QED ?}$$