

Comparing Marxen BusyBeaver with Ackermann function

Vladan Majerech

1.3.2010

Introduction

Remember definition of Ackerman function monotone in both variables:

$$A(x, 0) = 2$$

$$A(0, y) = y + 2$$

$$A(x + 1, y + 1) = A(x, A(x + 1, y))$$

One variable variant read from diagonal $A(x) = A(x, x)$. Remember that $A(1, y) = A(0, A(1, y - 1)) = A(1, y - 1) + 2 = A(1, 0) + 2y = 2y + 2 > 2y$, $A(2, y) = A(1, A(2, y - 1)) > 2A(2, y - 1) > 2^y A(2, 0) = 2^y \cdot 2 > 2^y$, $A(3, y) = A(2, A(3, y - 1)) > 2^{A(3, y - 1)} > (2 \uparrow y)(A(3, 0)) > 2 \uparrow y$ where $(2 \uparrow y)(x)$ denotes $f^{(y)}(x)$ for $f(x) = 2^x$ and $2 \uparrow y$ is shortcut for $f^{(y)}(1)$. $A(4, y) = A(3, A(4, y - 1)) > 2 \uparrow A(4, y - 1) > (2 \uparrow y)(A(4, 0)) > (2 \uparrow y)$ where $(2 \uparrow y)(x)$ denotes $f^{(y)}(x)$ for $f(x) = 2 \uparrow x$ and $2 \uparrow y$ is shortcut for $f^{(y)}(1)$. $A(6)$ is bigger than any number connected with the Universe.

Busy Beaver $S(n)$ function is the maximal number of steps deterministic Turing machine with n states working on $\{0, 1\}^*$ tape initialised with 0^* can process and stop. Our proof works also for function defined as the maximal number of ones deterministic Turing machine with n states working on $\{0, 1\}^*$ tape initialised with 0^* can left on tape and stop.

Sketch of proof that $S(61) > A(A(6))$

We will create TM with at most 61 states working on 0,1 tape initially filled with 0^* which will generate $0^*1^x0^*$ where $x > A(A(6))$. Obviously the running time would be higher.

Construction of TM will be divided into several parts. Original plan for first part was to prepare configuration $0^*10S \rightarrow 01^70^*$. Then computation of $A^{(i+1)}(w)$ would be invoked from state $0^*1^i0S \rightarrow 01^{w+1}0^*$. Computation of $A(w)$ transforms $|0S \rightarrow 01^{w+1}0^*$ to $|0^2b \rightarrow 1^{w+1}01^{w+1}0^*$. Third part should transform $|0^2b \rightarrow 1^{x+1}01^{y+1}0^*$ to $|\leftarrow z0^21^{A(x,y)+1}0^*$. Fourth part transforms $0^*1^i \leftarrow z0^21^{x+1}0^*$ to $0^*1^{i-1}0S \rightarrow 01^{x+1}0^*$ and stops for $|0 \leftarrow z0^21^{x+1}0^*$.

Alltogether $0^*I \rightarrow 0^* \Rightarrow 0^*10S \rightarrow 01^{6+1}0^* \Rightarrow 0^*10^2b \rightarrow 1^{6+1}01^{6+1}0^* \Rightarrow 0^*1 \leftarrow z0^21^{A(6,6)+1}0^* \Rightarrow 0^*0S \rightarrow 01^{A(6)+1}0^* \Rightarrow 0^*b \rightarrow 1^{A(6)+1}01^{A(6)+1}0^* \Rightarrow 0^* \leftarrow z0^21^{A(A(6),A(6))+1}0^* = 0^* \leftarrow z0^21^{A(A(6))+1}0^*$ where it stops.

Improvement to $S(35) > A(A(10))$

Finally I have reused some parts of program which allowed to reduce states count to 35. The first part using best Busy Beaver for 4 states prepares $0^*10S \rightarrow 11^{10+1}0^*$ instead. Actually I have resigned on construction of $A^{(i+1)}(w)$. I used $B^{(i+1)}(w)$ instead, where $B(w) = A(w, 2w + 3) > A(w)$. I have replaced second part with transformation $|0S \rightarrow (0|1)1^{w+1}0^*$ to $|0^2b \rightarrow 1^{w+1}01^{2w+3+1}0^*$. Moreover state pairs (b, d) can be unioned as well as (B, D) so there exists beaver with 35 states creating more than $A(A(6))$ ones (actually $B^{(2)}(10) > A(A(10))$).

See the commented transition function on the next page.

With 4 states Brady's initialisation we get $0^*I0^* \Rightarrow 0^*101^9 \leftarrow X1^30^* \Rightarrow 0^*10S \rightarrow 11^{10+1}0^* \Rightarrow 0^*10^2b \rightarrow 1^{10+1}01^{23+1}0^* \Rightarrow 0^*1 \leftarrow z0^21^{A(10,23)+1}0^* \Rightarrow 0^*S \rightarrow 01^{B(10)+1}0^* \Rightarrow 0^*b \rightarrow 1^{B(10)+1}01^{2B(10)+4}0^* \Rightarrow 0^* \leftarrow z0^21^{A(B(10),2B(10)+3)+1}0^* = 0^* \leftarrow z0^21^{B(B(10))+1}0^*$ being final configuration.

Marxen BusyBeaver $S(35) > Ackermann $A(A(10))$.$

state	tape 0	tape 1	comment	configuration
a	$0 \rightarrow b$	$1 \leftarrow a$	backing on first argument of $A(x,y)$	$ 01^{x+1-k} \leftarrow a 1^k 01^{y+1} 0^*$
			1 variable case (end)	$ 001^{w+1-k} \leftarrow a 1^k 01^{2w+4} 0^*$
$b = d$		$0 \rightarrow c$	stands on first 1 of first argument	$ 0b \rightarrow 1^{x+1} 01^{y+1} 0^*$
			when returned from r in nontrivial case	$ 01^x 0b \rightarrow 1^{x+1} 01^y 0^*$
			when returned from a on the same level	$ 0b \rightarrow 1^x 01^{A(x,y-1)+1} 0^*$
			1 variable case	$ 001^{w+1} 0b \rightarrow 1^{2w+2} 0^*$
			1 variable case (end)	$ 00b \rightarrow 1^{w+1} 01^{2w+4} 0^*$
c	$1 \leftarrow d$	$1 \rightarrow f$	was first argument 0?	$ 00c \rightarrow 1^x 01^{y+1} 0^*$
			1 variable case	$ 001^{w+1} 00c \rightarrow 1^{2w+1} 0^*$
$d = b$	$1 \leftarrow e$		case $A(0,y)=y+2$	$ 00 \leftarrow d 11^{y+1} 0^*$
e	$0 \leftarrow y$		finalizing $A(x,y)$	$ 0 \leftarrow e 1^{A(x,y)+1} 0^*$
f	$0 \rightarrow g$	$1 \rightarrow f$	running to y	$ 001^{x-k} f \rightarrow 1^k 01^{y+1} 0^*$
			1 variable case	$ 001^{w+1} 001^{2w+1-k} f \rightarrow 1^k 0^*$
g	$(1 \leftarrow D)$	$0 \rightarrow h$	0 finalizing for 1 variable else decrementing y	$ 001^x 0g \rightarrow 1^{y+1} 0^*$
			1 variable case (end)	$ 001^{w+1} 001^{2w+1} 0g \rightarrow 0^*$
h	$0 \leftarrow j$	$0 \rightarrow n$	was second argument 0?	$ 001^x 00h \rightarrow 1^y 0^*$
j	$0 \leftarrow j$	$0 \leftarrow k$	running to x	$ 001^x 0^* \leftarrow j 0^*$
k	$0 \rightarrow l$	$0 \leftarrow k$	clearing x	$ 001^* \leftarrow k 0^*$
l	$0 \rightarrow m$		clearing x	$ 00l \rightarrow 0^*$
m	$1 \leftarrow c$		case $A(x,0)=A(0,0)=2$	$ 000m \rightarrow 0^*$
n	$1 \leftarrow o$	$1 \rightarrow n$		$ 001^x 0001^{y-1-k} n \rightarrow 1^k 0^*$
o	$0 \leftarrow p$	$1 \leftarrow o$		$ 001^x 0001^{y-k} \leftarrow o 1^k 0^*$
p	$1 \leftarrow p$	$0 \leftarrow q$		$ 001^x 0^{2-k} \leftarrow p 1^k 01^y 0^*$
q	$1 \rightarrow r$	$1 \leftarrow q$	2 variables case	$ 01^j 01^{x-1-j-k} \leftarrow q 1^k 01^{2+j} 01^y 0^*$
			1 variable case	$ 001^j 01^{w-j-k} \leftarrow q 1^k 01^{2j+2} 0^*$
r	$0 \rightarrow b$	$0 \rightarrow s$	copy loop used both for 1 and 2 variables	$ 01^{j+1} r \rightarrow 1^{x-1-j} 01^{2+j} 01^y 0^*$
			at the end of (j) copy loop	$ 01^x r \rightarrow 01^{x+1} 01^y 0^*$
			1 variable case	$ 001^j r \rightarrow 1^{w-j+1} 01^{2j} 0^*$
s	$0 \rightarrow t$	$1 \rightarrow s$	2 variables case	$ 01^{j+1} 01^{x-2-j-k} s \rightarrow 1^k 01^{2+j} 01^y 0^*$
			1 variable case	$ 001^j 01^{w-j-k} s \rightarrow 1^k 01^{2j} 0^*$
t	$1 \rightarrow u$	$1 \rightarrow t$	2 variables case	$ 01^{j+1} 01^{x-2-j} 01^{2+j-k} t \rightarrow 1^k 01^y 0^*$
			1 variable case	$ 001^j 01^{w-j} 01^{2j-k} t \rightarrow 1^k 0^*$
u	$(1 \leftarrow x)$	$0 \rightarrow v$	0 for function with 1 variable, 1 for 2 variables	$ 01^{j+1} 01^{x-2-j} 01^{3+j} u \rightarrow 1^y 0^*$
			1 variable case	$ 001^j 01^{w-j} 01^{2j+1} u \rightarrow 0^*$
v	$1 \leftarrow w$	$1 \rightarrow v$		$ 01^{j+1} 01^{x-2-j} 01^{3+j} 01^{y-1-k} v \rightarrow 1^k 0^*$
w	$0 \leftarrow x$	$1 \leftarrow w$		$ 01^{j+1} 01^{x-2-j} 01^{3+j} 01^{y-k} \leftarrow w 1^k 0^*$
x	$0 \leftarrow q$	$1 \leftarrow x$	2 variables case	$ 01^{j+1} 01^{x-2-j} 01^{3+j-k} \leftarrow x 1^k 01^y 0^*$
			1 variable case	$ 001^j 01^{w-j} 01^{2j+2-k} \leftarrow x 1^k 0^*$
y	$0 \leftarrow z$	$1 \leftarrow a$	was this initial call?	$ 01^v \leftarrow y 01^{A(x,y)+1} 0^*$
z	$0 \rightarrow \star$	$1 \leftarrow A$	was this last $B(x)$ computation?	$0^* 1^{i-\ell} \leftarrow z 001^{B^{(\ell)}(w)+1} 0^*$
A	$0 \rightarrow B$	$1 \leftarrow A$	move through i	$0^* 1^{i-\ell-k} \leftarrow A 1^k 001^{B^{(\ell)}(w)+1} 0^*$
$B = D$		$0 \rightarrow C$	decrement i	$0^* B \rightarrow 1^{i-\ell} 001^{B^{(\ell)}(w)+1} 0^*$
C	$0 \rightarrow S$	$1 \rightarrow C$		$0^* 1^{i-\ell-1-k} C \rightarrow 1^k 001^{B^{(\ell)}(w)+1} 0^*$
S	$0 \rightarrow r$	$0 \rightarrow r$		$0^* 1^{i-\ell-1} 0S \rightarrow (0 1) 1^{B^{(\ell)}(w)+1} 0^*$
$D = B$	$1 \leftarrow E$		1 variable case (end)	$ 001^{w+1} 001^{2w+1} 0 \leftarrow D 10^*$
E	$1 \leftarrow F$	$1 \leftarrow E$	1 variable case (end)	$ 001^{w+1} 001^{2w+3-k} \leftarrow E 1^k 0^*$
F	$0 \leftarrow a$		1 variable case (end)	$ 001^{w+1} 0 \leftarrow F 1^{2w+4} 0^*$
X	$0 \rightarrow S$	$1 \leftarrow X$	Brady's record 4 states beaver	$0^* 101^{12-k} \leftarrow X 1^k 0^*$
I	$1 \rightarrow J$	$1 \leftarrow J$	Brady's A	
J	$1 \leftarrow I$	$0 \leftarrow K$	Brady's B	
K	$1 \rightarrow X$	$1 \leftarrow L$	Brady's C	
L	$1 \rightarrow L$	$0 \rightarrow I$	Brady's D	